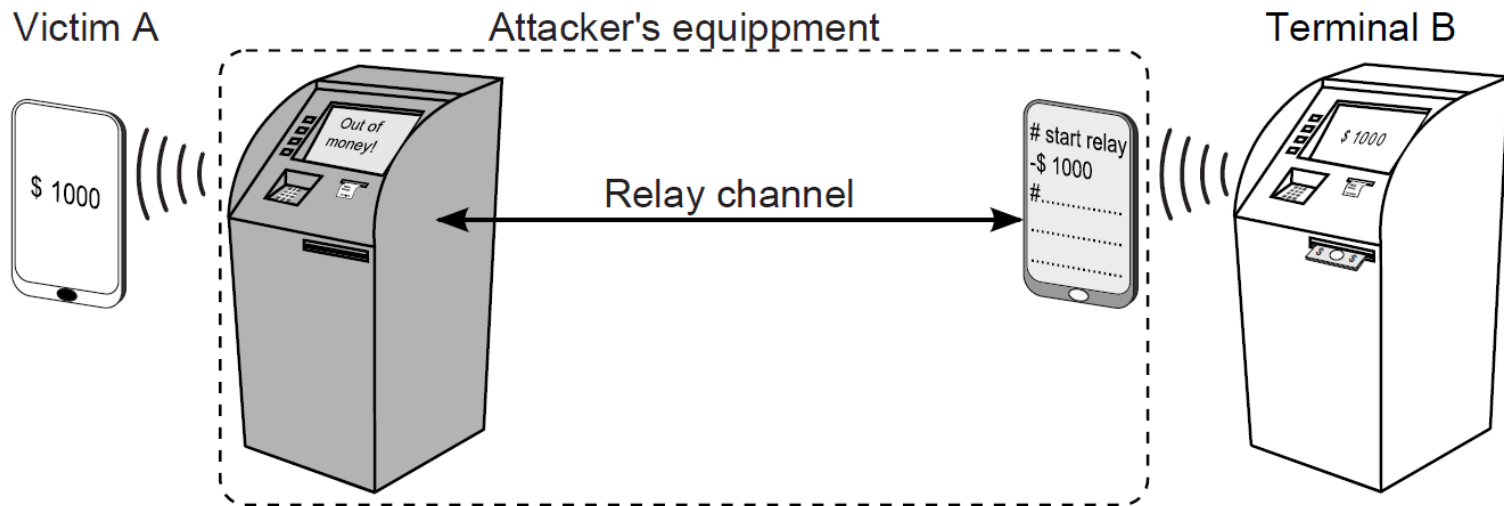


# Fortune Cookies and Smartphones



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Dani FESB-a, 2014

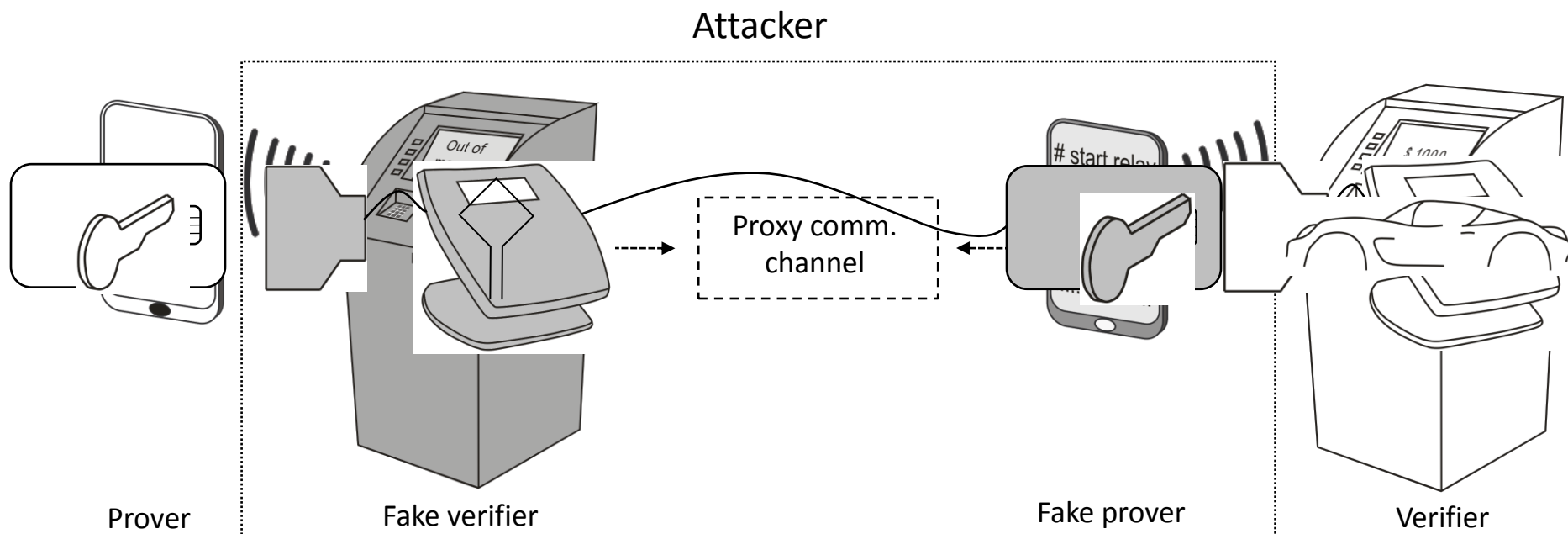
# Motivation

- Mobile payments
- Google + NFC = Google wallet
- AT&T + Verizon + T-Mobile = Isis
- NRC + QR codes



- La Caixa uses contactless terminals with NFC
- Zoosh uses ultrasound
- All these technologies are vulnerable to **mafia fraud attack**

# Motivation: Mafia-fraud attack



- Solutions to counter relay attacks
- RF distance bounding [6, 7]:
  - require modification of existing hardware
  - focused on a specific technology (e.g. RFID)
- Paradigm on unrelayed channels by Stajano et al. [8]
  - highly impracticable

[6] A. Francillon, B. Danev, S. Capkun. Relay Attacks on Passive Keyless Entry and Start Systems in Modern Cars. NDSS, 2011

[7] L. Francis, G. Hancke, K. Mayes, and K. Markantonakis. Practical NFC peer-to-peer relay attack using mobile phones. RFIDSec'10

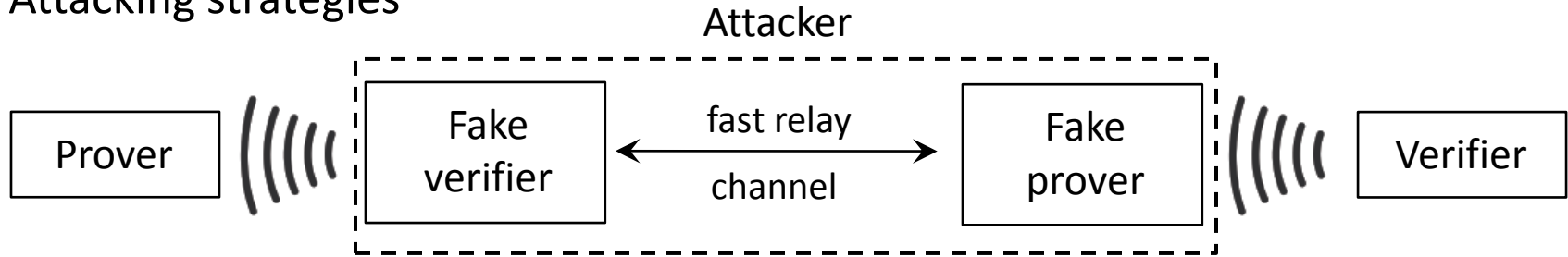
[8] F. Stajano, F.-L. Wong and B. Christianson, Multichannel Protocols to Prevent Relay Attacks, Financial Cryptography and Data Security, FC'10, 4–19, 2010.

# Desirable requirements

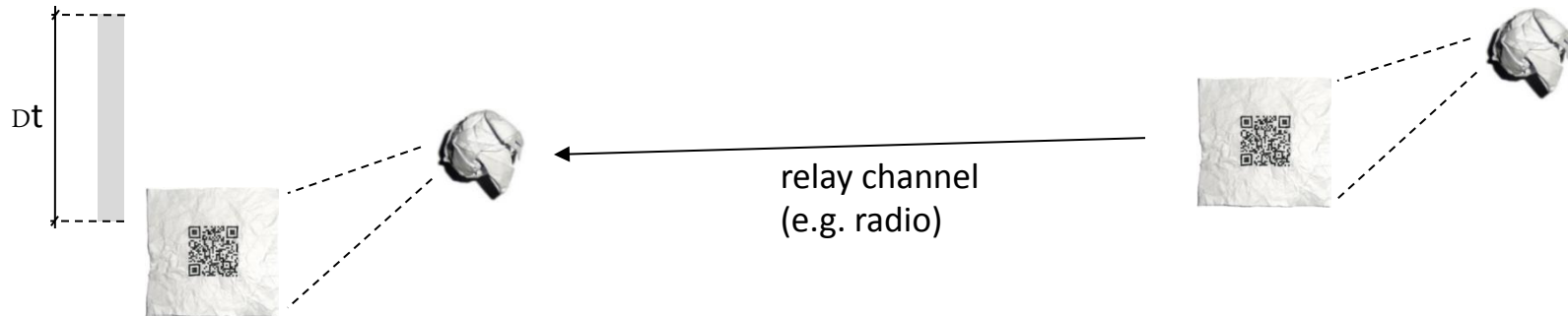
- Systems that use different communication technologies
- Minimal or no hardware changes to the existing equipment
- We use the paradigm on unrelayable channels based on multichannel protocols proposed by Stajano et al. [8]:
  - unclonability
  - unsimulability
  - untransportability
- We require the properties of unclonability and untransportability to hold for only a limited time period  $\Delta t$  – *weakly* unrelayable channel

# Weakly Unrelayable channel

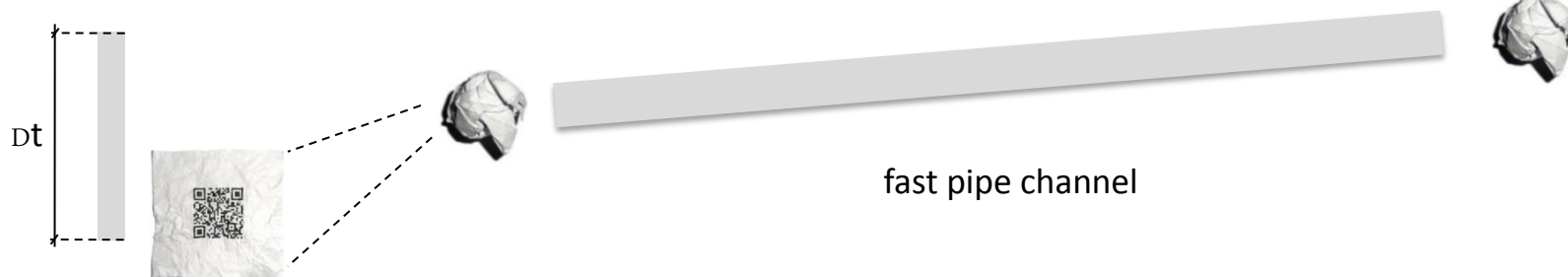
- Attacking strategies



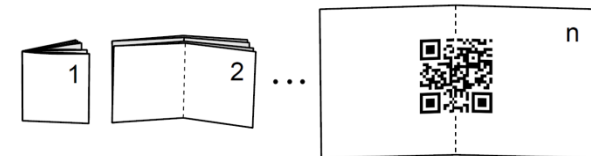
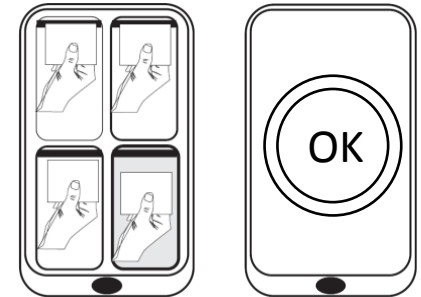
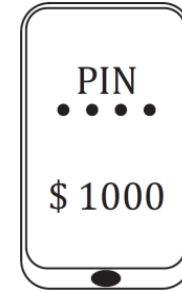
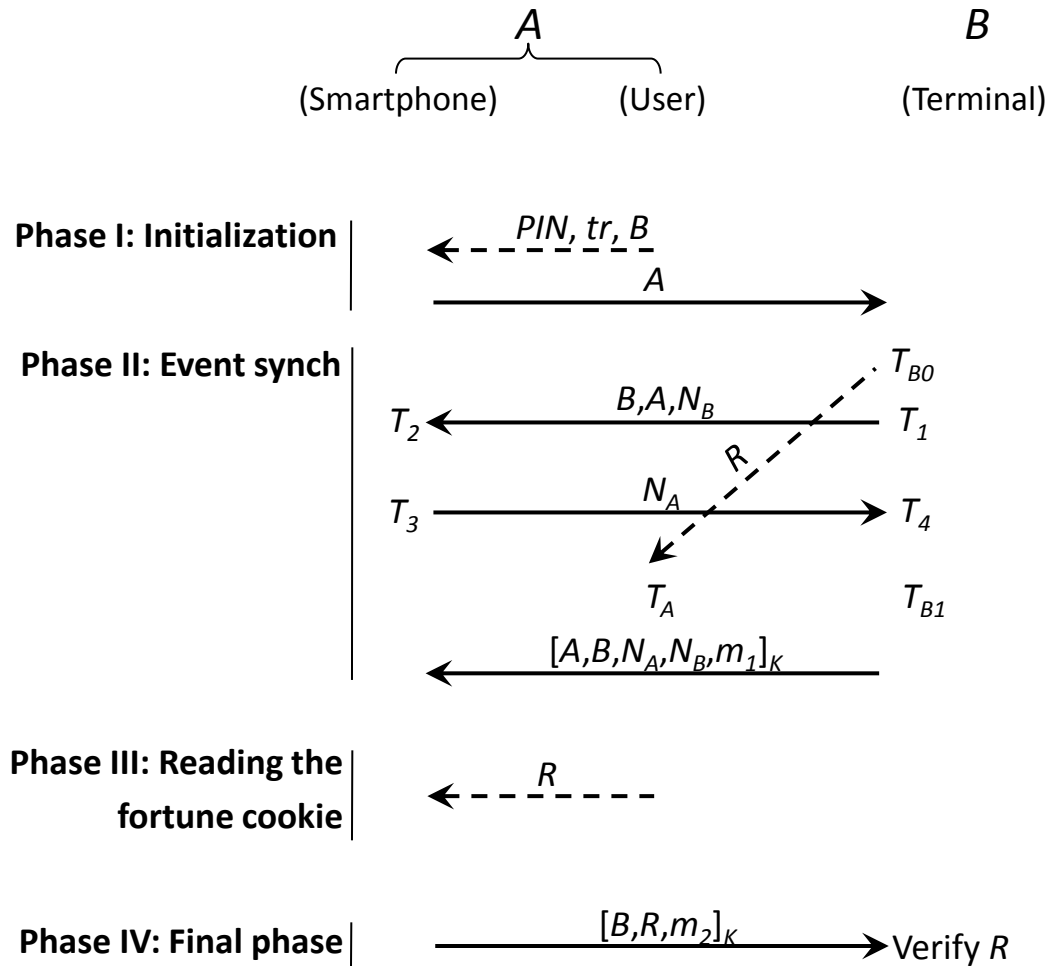
- Clone attack



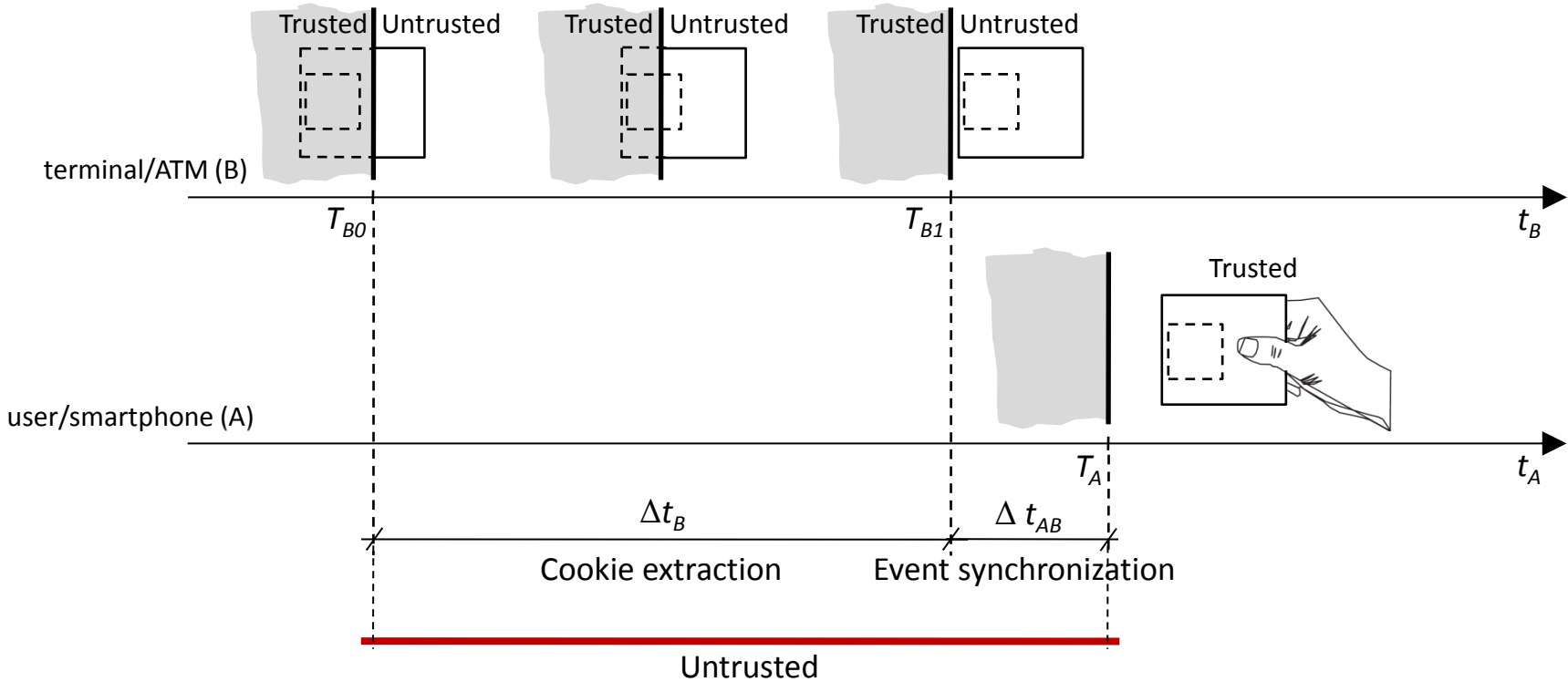
- Transport attack



# Authentication Protocol *Forces*

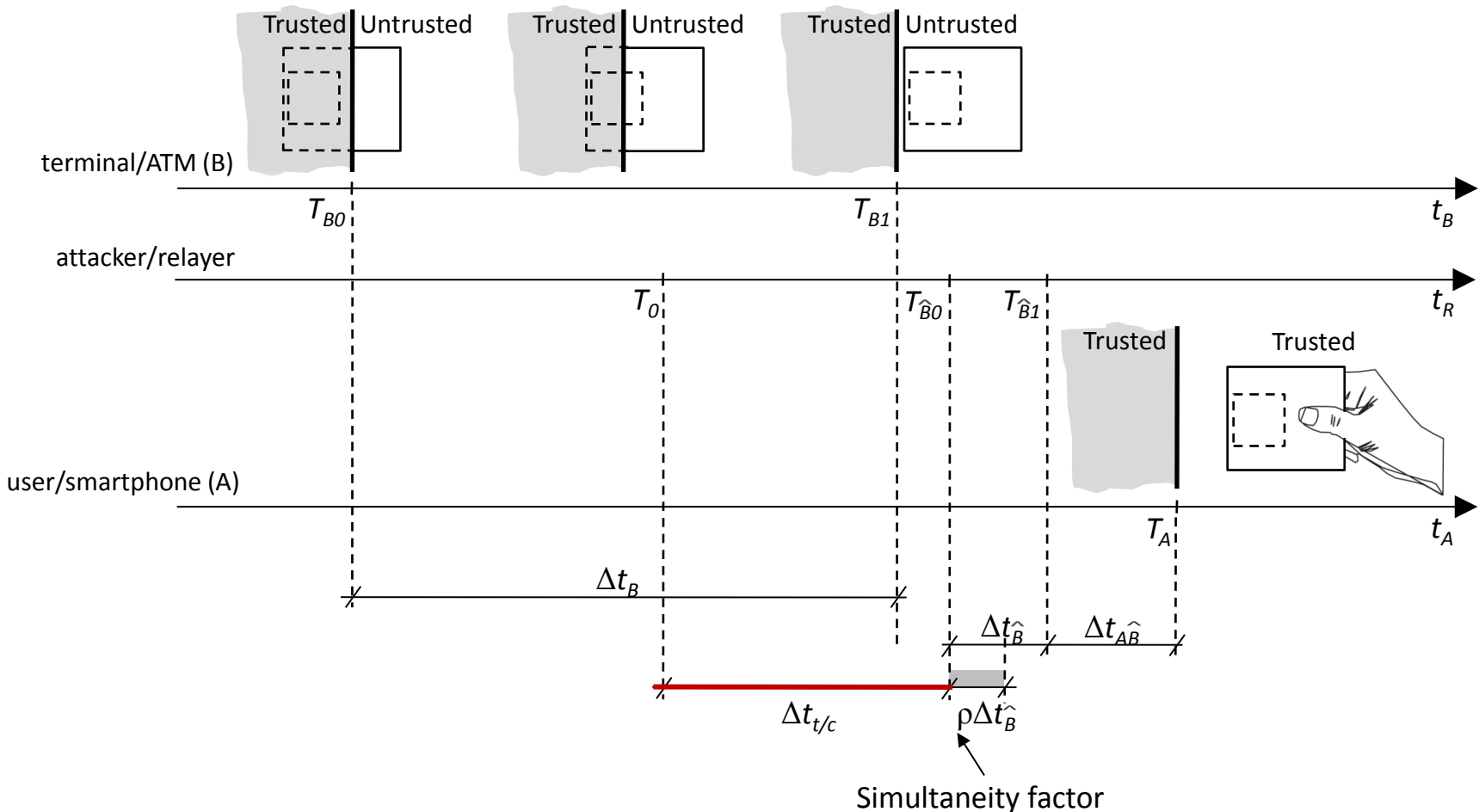


# Attacker Model



- Cookie extraction:  $\Delta t_B = T_{B1} - T_{B0} \leq \Delta t_B^*$
- Event synchronization:  $\Delta t_{AB} = T_A - T_{B1} \leq \Delta t_{AB}^*$

# Relay Attack Timing Model

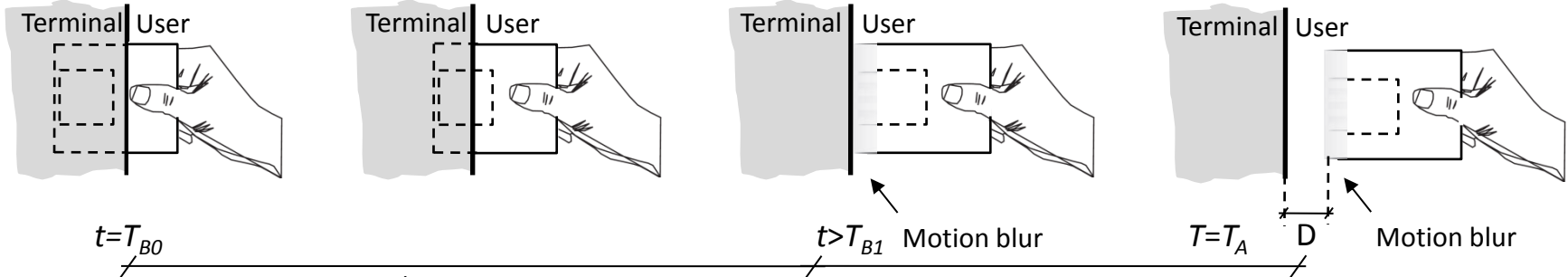


- the time window during which the attacker can mount the relaying attack:

$$\Delta t_{t/c} \leq \underline{(\Delta t_{AB}^* - \Delta t_{\hat{A}B})} + \underline{[\Delta t_B - (1-\rho) \Delta t_{\hat{B}}]} - (T_0 - T_{\hat{B}0})$$

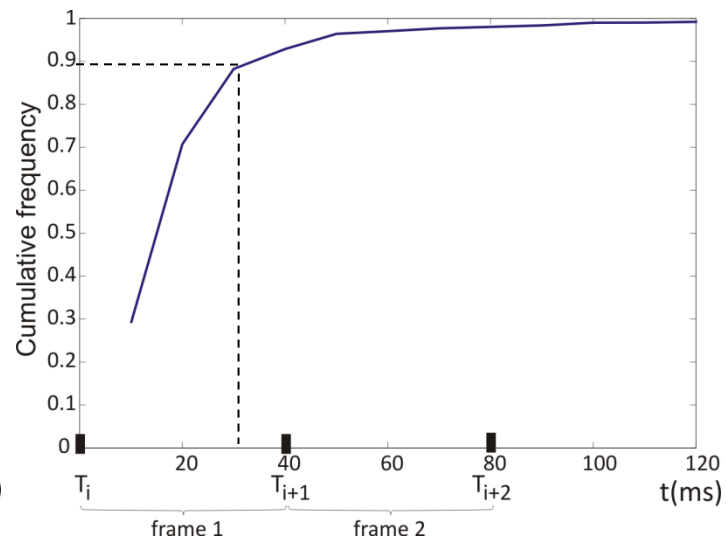
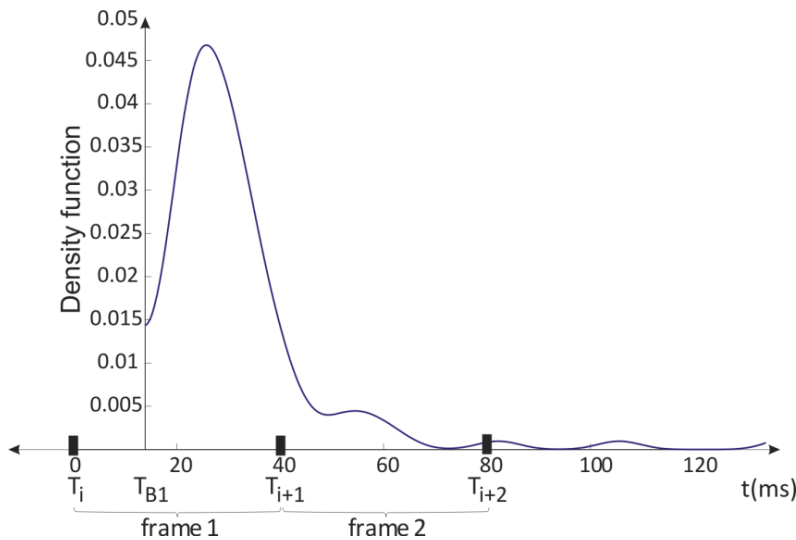


# Relay Attack Timing Model



$$\underline{\Delta t_B^*} = T_{B1} - T_{B0} \leq 120 \text{ ms}$$

$$\underline{\Delta t_{AB}^*} = \Delta t(D) + \Delta t_f \leq 70 \text{ ms}$$



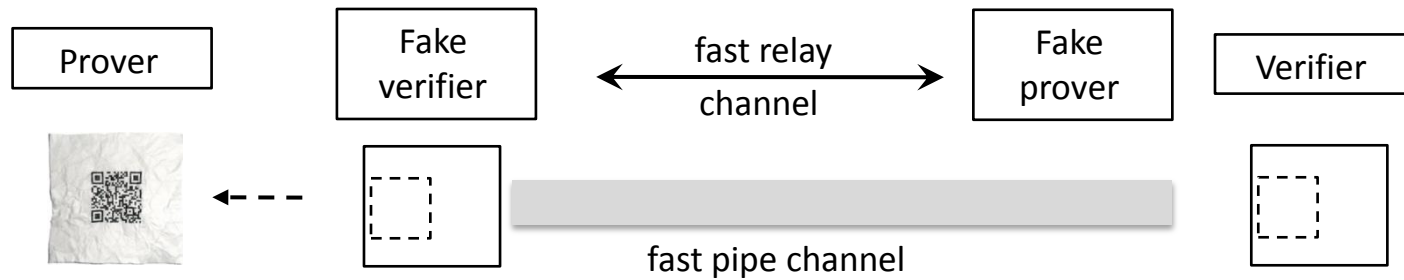
$$\Delta t_{AB} = T_A - T_{B1}$$

$$= \left( \left\lceil \frac{T_{B1} - T_i + \Delta t(D)}{\Delta t_f} - \alpha_D \right\rceil + \alpha_D \right) \Delta t_f + T_i - T_{B1}$$

$$\leq \left( \frac{T_{B1} - T_i + \Delta t(D)}{\Delta t_f} - \alpha_D + 1 + \alpha_D \right) \Delta t_f + T_i - T_{B1}$$

$$= \Delta t(D) + \Delta t_f .$$

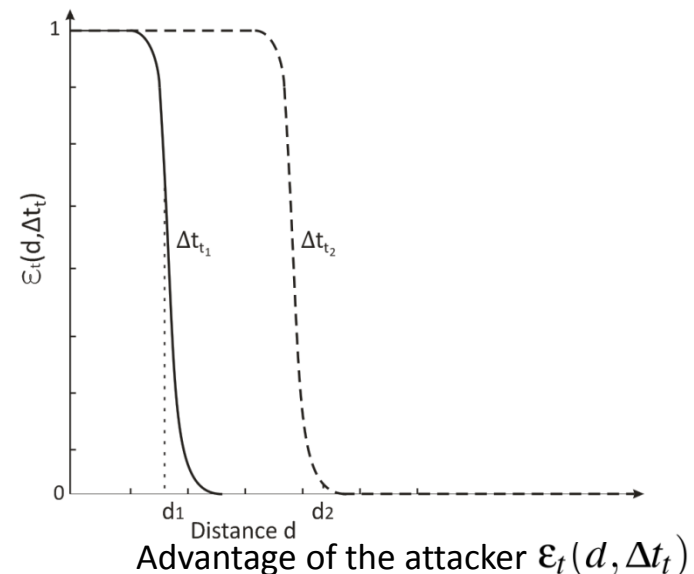
# Weak Untransportability



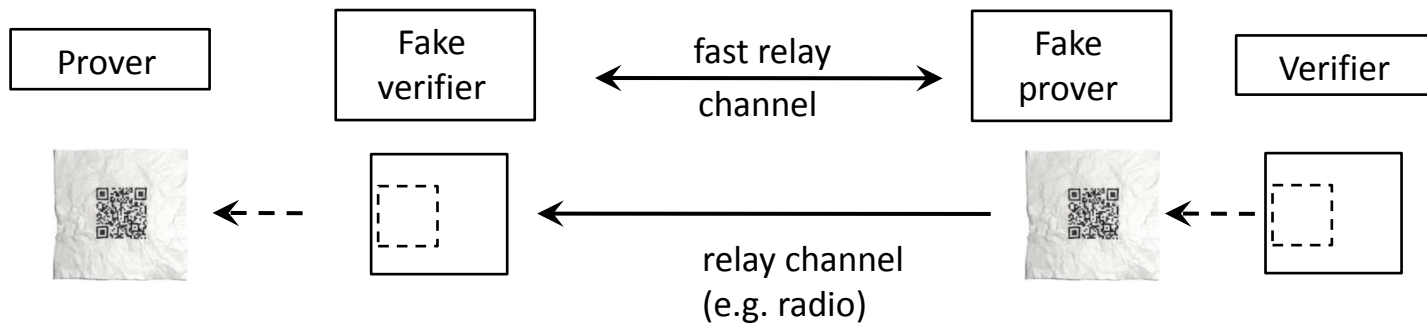
- Transport attack:  $\Delta t_t \leq (\Delta t_{AB}^* - \Delta t_B) = 70 \text{ ms} - 25 \text{ ms} = 45 \text{ ms}$
- Airbag inflates within approximately 60-80 ms
- Accessing an application to a secure element on a NFC-enabled phone takes 50-80 ms
- With “bang-bang” strategy the cookie can travel the maximum distance  $d$  within  $D\Delta t_t$

$$d(\Delta t_t) = \begin{cases} v_m(\Delta t_t - v_m/a) & \text{for } v_m < a\Delta t_t/2 \\ a \Delta t_t^2/4, & \text{otherwise} \end{cases}$$

$$a = 200 \text{ G} \quad v_m = 70 \text{ m/s} \quad d(\Delta t_t) = 1.01 \text{ m}$$

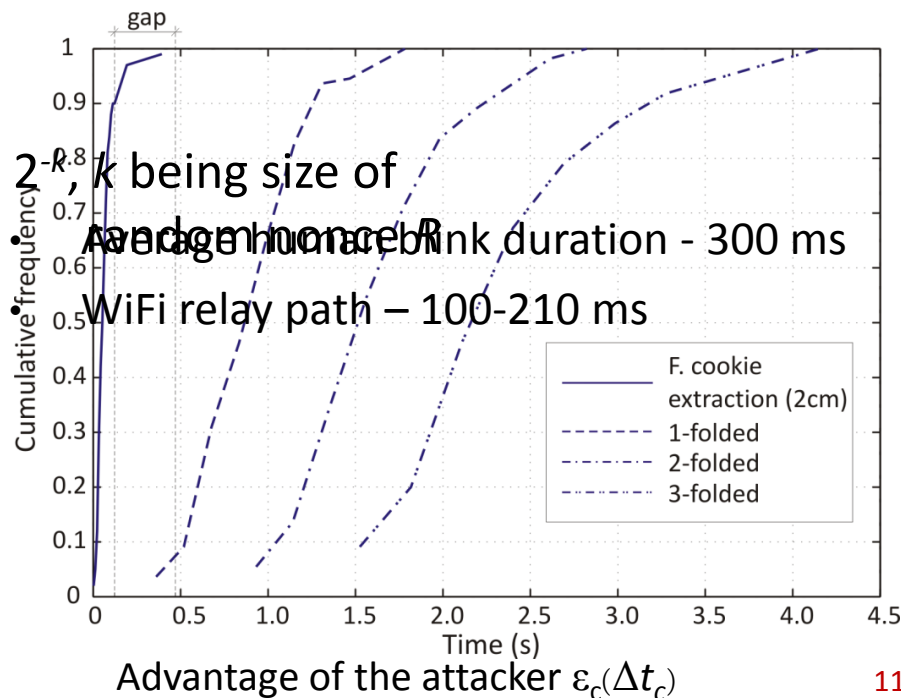


# Weak Uncloneability



- Clone attack:  $\Delta t_c \leq (\Delta t_{AB}^* - \Delta t_{AB}) + \Delta t_B^* = 70 \text{ ms} - 0 \text{ ms} + 120 \text{ ms} = 190 \text{ ms}$
- Timewise it is more optimal to launch the cloning attack (45ms vs 190 ms)

- Naive guessing attacker →
- Very strong X-Ray attacker →
- Realistic cloning attacker
- Powerful realistic cloning attacker



# Unsimulability



Fortune cookie



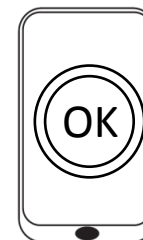
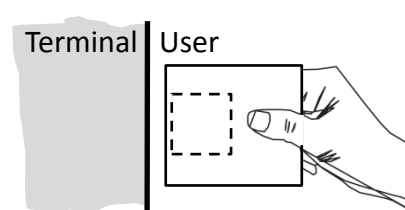
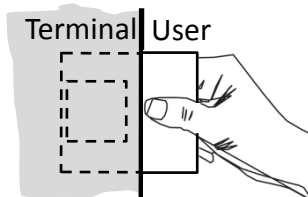
Smartphone



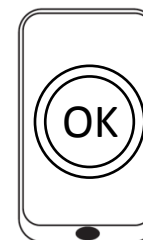
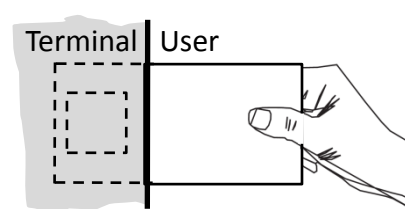
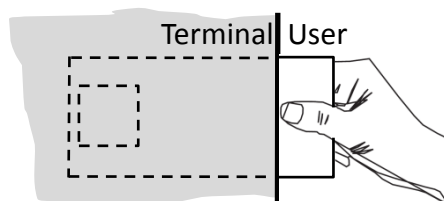
User

- A successful simulability attack buys the attacker an extra time
- Example: Longer fortune cookie size

- No attack



- Attack



- We managed to trick 25 percent of users
- Advantage of the attacker:  $\varepsilon_S(\varepsilon_t(d, \Delta t_t) + \varepsilon_c(\Delta t_c))$

# The Main Security Result

## Definition 1.

We say that a given protocol, executed between two honest parties  $A$  and  $B$ , is  $(d, \varepsilon)$  mafia fraud resistant, if the following holds except for the probability  $\varepsilon$ : a honest party  $B$  accepts to provide the service at time instant  $(t + \Delta t_s)$  **only if** (1) a honest party  $A$  (e.g., the user) has requested that service from  $B$  during period  $[t; t+\Delta t_s)$ , and (2)  $A$  and  $B$  have been within the distance  $d$  of each other during that period

- $\Delta t_s$  – allowable session duration
- $k$  – size of the random nonce
- $q$  – number of oracle calls
- $\varepsilon_t(\cdot)$  – probability of successful transport attack
- $\varepsilon_c(\cdot)$  – probability of successful cloning attack
- $\varepsilon_s(\cdot)$  – probability of successful simulability attack

## Theorem 1.

The Forces protocol is  $(d, \varepsilon)$  - mafia fraud resistant, with  $\varepsilon \leq q[(q+1)2^{-k} + \varepsilon_t(d, \Delta t_t) + \varepsilon_c(\Delta t_c) + \varepsilon_s(\varepsilon_t(d, \Delta t_{out}) + \varepsilon_c(\Delta t_{out}))]$

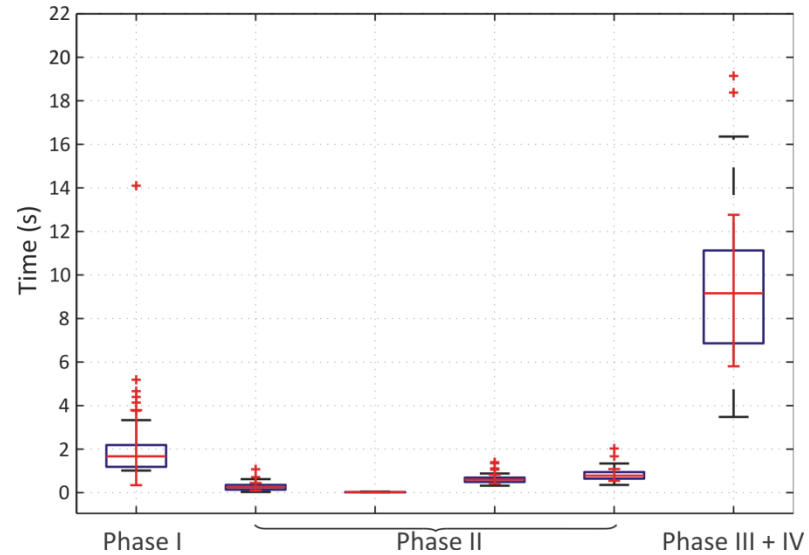
# User Performance Study

## Implementation

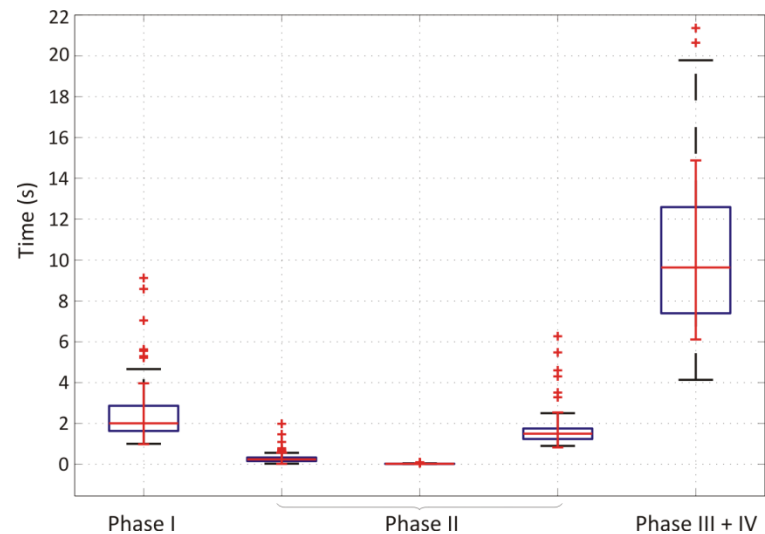


- Tests with 58 users
- Small execution times
- Average times around **13** seconds
- Small error rates

## Pushbutton-based event-synchronization



## Camera-based event-synchronization



# Conclusion

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- Design of a completely new protocol secure against relay attacks developed for financial transactions based on a paradigm of unrelayable channels,
- Formal proof of the protocol,
- Extensive user performance study on 54 users,
- The proposed solution has a reasonably low execution time (around 13 seconds on average), minimal error rate and is easy to understand,

**Thank you for your  
attention!**